## Lecture on Arithmetic Instructions (cont'd)

The SUB, SBB, DEC, AAS, DAS, NEG instructions
Subtraction instructions are similar to the addition instructions

| Mnemonic | Meaning | Format | Operation | Flags <br> affected |
| :---: | :---: | :---: | :---: | :---: |
| SUB | Subtract | SUB D,S | (D)-(S) $\longrightarrow(D)$ <br> Borrow $\longrightarrow(C F)$ | All |
| SBB | Subtract <br> with <br> borrow | SBB D,S | (D)-(S)-(CF) $\longrightarrow(D)$ | All |
| DEC | Decrement <br> by one | DEC D | (D)-1 $\longrightarrow(D)$ | All but CF |
| NEG | Negate | NEG D | All |  |
| DAS | Decimal <br> adjust for <br> subtraction | DAS | Convert the result in AL to <br> packed decimal format | All |
| AAS | ASCII <br> adjust for <br> subtraction | AAS | (AL) difference <br> (AH) dec by 1 if borrow | CY,AC |

The subtract (SUB) instruction is used to subtract the source from the destination The subtract with borrow (SBB) is similar to SUB but also subtract the carry flag From the destination.

EX: Assume the contents of registers BX and CX are 1234 H and 0123 H , respectively, and the carry flag is 0 , what is the result of executing the following instruction SBB BX, CX

The instruction implements the operation

$$
(B X)-(C X)-(C F) \quad \longrightarrow(B X)
$$

Therefore, $(B X)=1234 \mathrm{H}-0123 \mathrm{H}-0=1111 \mathrm{H}$
The carry flag remains cleared since no borrow is needed.

The decrement (DEC) instruction is used to subtract 1 from its operand. It does not affect the carry flag.

EX: Assume BX has 0000 H, What is the results of executing the following instruction: DEC BX

Subtracting 1 from 0000H would produce FFFFH However, the carry flag will remain unchanged.

The negate (NEG) instruction replaced its operand by its negative.
The carry flag will be become 1 as a result of this instruction.

- Assume $(B X)=003 A H$, what is the result of executing the instruction

NEG BX
$(B X)=0000000000111010$

$$
\begin{aligned}
2 ' s \text { comp } & =1111111111000110= \\
& =\text { FFC } 6 \mathrm{H} \text { and }(\mathrm{CF})=1
\end{aligned}
$$

The SUB and SBB instruction can subtract numbers represented in ASCII and BCD. Just like in addition, the results must be adjusted To produce the corresponding decimal numbers. The instructions AAS, and DAS are provided to perform the adjustments in AL.

EX: MOV BL, 28H<br>MOV AL, 83H<br>SUB AL,BL; AL=5BH<br>DAS ; adjust as $\mathrm{AL}=55 \mathrm{H}$

EX: MOV AX, 38H
SUB AL,39H; AX=00FF
AAS ; AX=FF09 ten's complement of -1 (Borrow one from AH )
OR AL, 30H ; AL=39

EX: 32-bit subtraction of two 32 bit numbers $X$ and $Y$ that are stored in the memory as

$$
\begin{aligned}
& X=(D S: 203 h)(D S: 202 h)(D S: 201 h)(D S: 200 h) \\
& Y=(D S: 103 h)(D S: 102 h)(D S: 101 h)(D S: 100 h)
\end{aligned}
$$

- The result $X-Y$ to be stored where $X$ is saved in the memory

```
    MOV SI, 200h
    MOV DI, 100h
    MOV AX, [SI]
    SUB AX, [DI]
    MOV [SI], AX ;save the LS word of result
    MOV AX, [SI] +2
    SBB AX, [DI]+2
    MOV [SI] +2, AX
Ex. 123456 78-2345 67 89 = EF EE EE EE
```


## Multiplication and Division

Multiplication and division can be performed on signed or unsigned numbers For unsigned numbers, MUL and DIV instructions are used, while for signed numbers IMUL and IDIV are used.

The format of the multiplication and division instructions does not specify the multiplicand as it is implicitly specified depending on the size of the source As shown in the following tables.

| Multiplication <br> (MUL or IMUL) | Multiplicand | Operand <br> (Multiplier) | Result |
| :---: | :---: | :---: | :---: |
| Byte*Byte | AL | Register or <br> memory | AX |
| Word*Word | AX | Register or <br> memory | DX :AX |


| Division <br> (DIV or IDIV) | Dividend | Operand <br> (Divisor) | Quotient: <br> Remainder |
| :---: | :---: | :---: | :---: |
| Word/Byte | AX | Register or <br> Memory | AL: AH |
| Dword/Word | DX:AX | Register or <br> Memory | AX : DX |

For the unsigned division, if the quotient is larger than FF in the division By a byte and FFFF in the division by a word type 0 interrupt occurs.

For the signed division, if the quotient is positive and exceeds 7 F for the byte division or 7FFF for the word division type 0 interrupts occurs. However, if the quotient is negative and less than 81 H for byte division and 8001 H for word division, then type 0 interrupt occurs.

The following examples explains the steps to perform multiplication and division Using calculators to give same results as the 88/86 processors.

Ex1: Assume that each instruction starts from these values:
$\mathrm{AL}=85 \mathrm{H}, \mathrm{BL}=35 \mathrm{H}, \mathrm{AH}=0 \mathrm{H}$

1. $\mathrm{MUL} \mathrm{BL} \rightarrow \mathrm{AL} . \mathrm{BL}=85 \mathrm{H} * 35 \mathrm{H}=1 \mathrm{~B} 89 \mathrm{H} \rightarrow \mathrm{AX}=1 \mathrm{~B} 89 \mathrm{H}$
2. IMUL BL $\rightarrow \mathrm{AL} . \mathrm{BL}=2^{\prime} \mathrm{S}$ AL * $\mathrm{BL}=2^{\prime} \mathrm{S}(85 \mathrm{H}) * 35 \mathrm{H}$ $=7 \mathrm{BH} * 35 \mathrm{H}=1977 \mathrm{H} \rightarrow 2^{\prime}$ s comp $\rightarrow \mathrm{E} 689 \mathrm{H} \rightarrow \mathrm{AX}$.
3. $\mathrm{DIV} \mathrm{BL} \rightarrow \frac{A X}{B L}=\frac{0085 H}{35 H}=02(85-02 * 35=1 \mathrm{~B}) \rightarrow$| AH | AL |
| :--- | :--- |
| 1 B | 02 |
4. IDIV $\mathrm{BL} \rightarrow \frac{A X}{B L}=\frac{0085 H}{35 H}=$| AH | AL |
| :--- | :--- |
| 1 B | 02 |

Ex2: $\quad A L=F 3 H, B L=91 H, A H=00 H$

1. $\mathrm{MUL} \mathrm{BL} \rightarrow \mathrm{AL} * \mathrm{BL}=\mathrm{F} 3 \mathrm{H} * 91 \mathrm{H}=89 \mathrm{~A} 3 \mathrm{H} \rightarrow \mathrm{AX}=89 \mathrm{~A} 3 \mathrm{H}$
2. $\mathrm{IMUL} \mathrm{BL} \rightarrow \mathrm{AL} * \mathrm{BL}=2^{\prime} \mathrm{S} A L * 2^{\prime} \mathrm{S} \mathrm{BL}=2^{\prime} \mathrm{S}(\mathrm{F} 3 \mathrm{H}) * 2^{\prime} \mathrm{S}(91 \mathrm{H})=$ $0 \mathrm{DH} * 6 \mathrm{FH}=05 \mathrm{~A} 3 \mathrm{H} \rightarrow \mathrm{AX}$.
3.IDIV BL $\rightarrow \frac{A X}{B L}=\frac{00 F 3 H}{2^{\prime} S(91 H)}=\frac{00 F 3 H}{6 F H}=2 \rightarrow(00 \mathrm{~F} 3-2 * 6 \mathrm{~F}=15 \mathrm{H})$

| AH | AL |
| :---: | :---: |
| 15 | 02 |
| $\mathbf{R}$ | $\mathbf{Q}$ |$\rightarrow \frac{P O S}{N E G}=N E G \rightarrow 2^{\prime} \mathrm{s}(02)=\mathrm{FEH} \rightarrow$| AH | AL |
| :---: | :---: |
| 15 | FE |

4. DIV BL $\rightarrow \frac{A X}{B L}=\frac{00 F 3 H}{91 H}=01 \rightarrow(\mathrm{~F} 3-1 * 91=62) \rightarrow$

| AH | AL |
| :---: | :---: |
| 62 | 01 |
| $R$ | $Q$ |

## Ex3: $\quad A X=F 000 H, B X=9015 H, D X=0000 H$

1. $\mathrm{MUL} \mathrm{BX}=\mathrm{F} 000 \mathrm{H} * 9015 \mathrm{H}=8713 \mathrm{~B}=00$
2. $\operatorname{IMUL} \mathrm{BX}=2^{\prime} \mathrm{S}(\mathrm{F} 000 \mathrm{H}) * 2^{\prime} \mathrm{S}(9015 \mathrm{H})=1000 * 6 \mathrm{FEB}=$ 06FE B000
3. DIV BL $=\frac{F 000 H}{15 H}=\mathrm{B} 6 \mathrm{DH} \rightarrow$ More than $\mathrm{FFH} \rightarrow$ Divide Error.
4. IDIV BL $\rightarrow \frac{2^{\prime} S(F 000 H)}{15 H}=\frac{1000 H}{15 H}=\mathrm{C} 3 \mathrm{H}>7 \mathrm{~F} \rightarrow$ Divide Error.

## Ex4: AX=1250H, BL=90H

$$
\text { 1. IDIV } \begin{aligned}
\mathrm{BL} \rightarrow \frac{A X}{B L} & =\frac{1250 H}{90 H}=\frac{P O S}{N E G}=\frac{P O S}{2^{\prime} s N E G}=\frac{1250 H}{2^{\prime} s(90 H)}=\frac{1250 H}{70 H} \\
& =29 \mathrm{H}(\mathrm{Q}) \rightarrow(1250-29 * 70)=60 \mathrm{H}(\mathrm{REM})
\end{aligned}
$$

$$
29 \mathrm{H}(P O S) \rightarrow 2^{\prime} \mathrm{S}(29 \mathrm{H})=\mathrm{D} 7 \mathrm{H} \rightarrow \begin{array}{|cc}
\mathrm{R} & \mathrm{Q} \\
\hline 60 \mathrm{H} & \mathrm{D} 7 \mathrm{H}
\end{array}
$$

2. DIV $\mathrm{BL} \rightarrow \frac{A X}{B L}=\frac{1250 \mathrm{H}}{90 H}=20 \mathrm{H} \rightarrow 1250-20^{*} 90=50 \mathrm{H} \rightarrow$| R | Q |
| :---: | :---: |
| 50 H | 20 H |
| AH | AL |

To divide an 8-bit dividend by and 8-bit divisor by extending the sign bit Of Al to fill all bits of AH . This can be done automatically by executing the Instruction (CBW).

In a similar way 16 -bit dividend in $A X$ can be divided by 16-bit divisor. In this case the sign bit in AX is extended to fill all bits of DX. The instruction CWD perform this operation automatically.

Note that CBW extend 8 -bit in $A L$ to 16-bit in $A X$ while the value in $A X$ will Be equivalent to the value in AL. Similarly, CWD convert the value in $A X$ to 32-bit In ( $D X, A X$ ) without changing the original value.

